

Name: _____ Email: _____

Problem 1 (18%). Assemble the following machine instructions into **binary**, use spaces to separate fields, and **registers** in their symbolic form (\$ra NOT \$31). Assume absolute jump addresses.

Field 1	Fields 2 and etc	instruction
001011 (0xb)	11101 00001 0000 0000 0000 0011 (\$sp) (\$at) (3)	sltiu \$at,\$sp,3
000000 (0)	0000000000 11111 00000 010000 (0) (\$ra) (0) (0x10)	mfhi \$ra
000000 (0)	00000 01001 10001 00000 100010 (\$0) (\$t1) (\$s1) (0) (0x22)	neg \$s1,\$t1 (2's complement s1= -t1;) sub \$s1,\$0,\$t1
000010 (2)	0000000000 0001 1110 1011 1111 (0x1ebf) = word offset	j 0x7afc (0x7afc = byte offset)
101001 (0x29)	11110 00010 0000 0000 0000 1100 (\$fp) (\$v0) (12)	sh \$v0,12(\$fp)
0x03e00008 (0) (\$ra) (0) (8)	000000 11111 0000 0000 0000 0000 01000 (0) (\$ra) (0) (8)	jr \$ra

Problem 2 (7%).

Assume each part is **independent**. Assume absolute jump & branch addresses (no pc relative).

Fill in only registers that changed!

What is the value of the register or memory contents **after** the execution of the instruction.

Assume pc = 200; \$s3=13; \$s4=6; \$ra=250; memory[8]=0xfedcba98; memory[12]=0x76543210;

Instruction	pc	\$ra	\$s3	\$s4	memory[8]	Memory[12]
srl \$s3, \$s4, 3	204		0			
lui \$s3, 0xffff	204		0xFFFF0000			
lh \$s4, 8(\$s4)	204			0x3210		
addi \$s3, \$s3,-4	204		9			
sb \$s4, 7(\$s4)	204					0x76063210
bnez \$ra, 600	600					
j 250	250					

Problem 3. (25%) Translate the following C code into MIPS. Please comment your code. Assume signed unless defined otherwise. The value **x** is \$s0; the value **y** is \$s1; the value **s** is \$s2; **t** is \$s3; the address of **r** is in \$s4; **p** is \$s5; the address of **d** is in \$s6, **w** is \$s7.

Points will be taken off for assembler syntax errors.

register unsigned int x, y; register int s, *t; struct { int a[3]; short b; } r, *p; short d[5]; char *w;

(a) `s = *t;`

```
lw $s2, 0($s3)    #s= (int) *t
```

(b) `s = *w;`

```
lb $s2, 0($s7)    #s = (char) *w
```

(c) `s += d[x + 3];`

```
addiu $t0, $s0, 3    #t0=x+3
addu $t1, $t0, $t0   #t1=(x+3)*sizeof(short) = (x+3)*2
addu $t3, $s6, $t1   #t3={address of d+(x+3)*sizeof(short)}
lh $t4, 0($t3)       #t4= (short) d[x+3]
add $s2, $s2, $t4    #s = (signed int) {s + d[x+3]}
```

(d) `*(d + 3) = s + r.b; /* offsets within struct 0 : a[0], 4: a[1], 8:a[2], 12 : b */`

```
lh $t0, 12($s4)     #t0= (short) r.b
add $t1, $s2, $t0   #t1= (signed int) {s+r.b}
sh $t1, 6($s6)      #*(d+3)= (short)d[3]{6=3*sizeof(short)}
```

(e) `for(x=0; x <= 10; x++) { s = (s >> 13); } /* or {s >>= 13} */`

```
addi $t0, $0, 10    #t0 = 10
addu $s0, $0, $0    #(unsigned int ) x =0
L2: bgt $s0, $t0,L1  #if ( x > 10) goto L1
sra $s2, $s2,13     #(signed int)s >> 13
addiu $s0, $s0, 1   #(unsigned int) x++;
j L2
L1:
```

Problem 4. (25%) Translate the following code and add **comments**
Points will be taken off for assembler syntax errors.

```
int  unicodestrln(short *s) {
    register int  n = 0;

    for(i=0;  s[i] != 0;  i++) {n++; }

    return n;
}
```

(a) Write the prolog

#prolog is empty because:

- \$t0..\$t4 registers are only used (by convention these registers are not required to be saved).
- Function does not call another function, therefore no need to save \$ra.

(b) Write the body (hint: write the body first; then write the prolog)

```

    addi $t0,$0,0      #n=0;
    addi $t1,$0,$0     #i=0;
L2:  sll  $t2,$t1,1     #t2 = i * sizeof(short)
    add  $t3,$a0,$t2   #t3 = address(s + i)
    lh  $t4,0($t3)    #t4 = *(s+i) = s[i]
    beq $t4,$0,L1
    addi $t0,$t0,1     #n++;
    addi $t1,$t1,1     #i++;
    j   L2
L1:  addi $v0,$t0,0    #return n

```

add \$t2,\$t1,\$t1

(c) Write the epilog

```
END: jr $ra          #return
```

Alternative Solution

```

    addi $v0, $0, 0
L2:  lh    $t4, 0($a0)
    beq   $t4, $0, END
    addi  $v0, $0, 1
    addi  $a0, $0, 2 /*2 = sizeof(short)*/
    j     L2

END: jr $ra

```

Problem 5. (10%) Translate the following global variables and assign the location counter beginning at 700

```
(a) class xxstring {
    char    argv;
    short   montab[2];
    short   (*daytab)[13];
    void    (*strcpy)();
    class   xxstring **next;
} *fsp;
```

Decimal Location	Assembler definitions
700	<code>fsp : .word 0</code>
704	

```
(b) class ccstring {
    char    argv;
    short   montab[2];
    short   (*daytab)[13];
    void    (*strcpy)();
    class   ccstring **next;
} *fcp;
```

Decimal Location	Assembler definitions
700	<code>fcp : .word 0</code>
704	

Problem 6. (15%) Given the following instruction sequence in the table below.

Assume the (alu and slt instructions are 5 clocks); (loads 10 clocks); (stores 20 clocks); (jumps 2 clocks); (branches 3 fall through/6 for branch);

- Show the **best** case timing path through the code showing annotations and total.
- Show the **worst** case timing path through the code showing annotations and total.
- What values for \$s1,\$s2,\$s3,\$s4 will make this code execute the worst case?

`$s1 < $s2 ; $s3,$s4 any values`

Instruction	best case	worst case
<code>slt \$t1,\$s1,\$s2</code>	5	5
<code>bne \$t1,\$0,L1</code>	3	6
<code>addi \$t2,\$zero,5</code>	5	X
<code>j L2</code>	2	X
L1: <code>slt \$t2,\$s3,\$s4</code>	X	5
<code>beq \$0,\$0,L2</code>	X	6
<code>xori \$s3,\$s4,3</code>	X	X
L2: <code>addi \$s1,\$zero,10</code>	5	5
Total Time	20	27